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1. Teoría de números

1.1. Big mod

```
//retorna (b^p)mod(m)  
// 0 <= b,p <= 2147483647  
// 1 <= m <= 46340  
long f(long b, long p, long m){  
  long mask = 1;  
  long pow2 = b % m;  
  long r = 1;  
  
  while (mask){  
    if (p & mask)  
      r = (r * pow2) % m;  
    pow2 = (pow2*pow2) % m;  
    mask <<= 1;  
  }  
  return r;  
}
```

1.2. Criba de Eratóstenes

Marca los números primos en un arreglo. Algunos tiempos de ejecución:

SIZE	Tiempo (s)
100000	0.004
1000000	0.078
10000000	1.550
100000000	14.319

```
#include <iostream>

const int SIZE = 1000000;

//criba[i] = false si i es primo
bool criba[SIZE+1];

void buildCriba(){
    memset(criba, false, sizeof(criba));

    criba[0] = criba[1] = true;
    for (int i=2; i<=SIZE; i += 2){
        criba[i] = true;
    }

    for (int i=3; i<=SIZE; i += 2){
        if (!criba[i]){
            for (int j=i+i; j<=SIZE; j += i){
                criba[j] = true;
            }
        }
    }
}
```

1.3. Divisores de un número

Este algoritmo imprime todos los divisores de un número (en desorden) en $O(\sqrt{n})$. Hasta 4294967295 (máximo *unsigned long*) responde instantaneamente. Se puede forzar un poco más usando *unsigned long long* pero más allá de 10^{12} empieza a responder muy lento.

```
for (int i=1; i*i<=n; i++) {
    if (n%i == 0) {
        cout << i << endl;
        if (i*i<n) cout << (n/i) << endl;
    }
}
```

2. Grafos

2.1. Algoritmo de Dijkstra

El peso de todas las aristas debe ser no negativo.

```
#include <iostream>
#include <algorithm>
#include <queue>

using namespace std;
```

```

struct edge{
    int to, weight;
    edge() {}
    edge(int t, int w) : to(t), weight(w) {}
    bool operator < (const edge &that) const {
        return weight > that.weight;
    }
};

int main(){
    int N, C=0;
    scanf("%d", &N);
    while (N-- && ++C){
        int n, m, s, t;
        scanf("%d%d%d%d", &n, &m, &s, &t);
        vector<edge> g[n];
        while (m--){
            int u, v, w;
            scanf("%d%d%d", &u, &v, &w);
            g[u].push_back(edge(v, w));
            g[v].push_back(edge(u, w));
        }

        int d[n];
        for (int i=0; i<n; ++i) d[i] = INT_MAX;
        d[s] = 0;
        priority_queue<edge> q;
        q.push(edge(s, 0));
        while (q.empty() == false){
            int node = q.top().to;
            int dist = q.top().weight;
            q.pop();

            if (dist > d[node]) continue;
            if (node == t) break;

            for (int i=0; i<g[node].size(); ++i){
                int to = g[node][i].to;
                int w_extra = g[node][i].weight;

                if (dist + w_extra < d[to]){
                    d[to] = dist + w_extra;
                    q.push(edge(to, d[to]));
                }
            }
        }
        printf("Case #%d: ", C);
        if (d[t] < INT_MAX) printf("%d\n", d[t]);
        else printf("unreachable\n");
    }
    return 0;
}

```

2.2. Minimum spanning tree: Algoritmo de Prim

```

#include <stdio.h>
#include <string>

```

```

#include <set>
#include <vector>
#include <queue>
#include <iostream>
#include <map>

using namespace std;

typedef string node;
typedef pair<double, node> edge;
typedef map<node, vector<edge> > graph;

int main(){
    double length;
    while (cin >> length){
        int cities;
        cin >> cities;
        graph g;
        for (int i=0; i<cities; ++i){
            string s;
            cin >> s;
            g[s] = vector<edge>();
        }
        int edges;
        cin >> edges;
        for (int i=0; i<edges; ++i){
            string u, v;
            double w;
            cin >> u >> v >> w;
            g[u].push_back(edge(w, v));
            g[v].push_back(edge(w, u));
        }

        double total = 0.0;
        priority_queue<edge, vector<edge>, greater<edge> > q;
        q.push(edge(0.0, g.begin()->first));
        set<node> visited;
        while (q.size()){
            node u = q.top().second;
            double w = q.top().first;
            q.pop(); ///

            if (visited.count(u)) continue;

            visited.insert(u);
            total += w;
            vector<edge> &vecinos = g[u];
            for (int i=0; i<vecinos.size(); ++i){
                node v = vecinos[i].second;
                double w_extra = vecinos[i].first;
                if (visited.count(v) == 0){
                    q.push(edge(w_extra, v));
                }
            }
        }
    }
}

```

```

    if (total > length){
        cout << "Not enough cable" << endl;
    }else{
        printf("Need%.11f miles of cable\n", total);
    }
}
return 0;
}

```

2.3. Minimum spanning tree: Algoritmo de Kruskal + Union-Find

```

#include <iostream>
#include <vector>
#include <algorithm>

using namespace std;

/*
Algoritmo de Kruskal, para encontrar el árbol de recubrimiento de mínima suma.
*/

struct edge{
    int start, end, weight;
    bool operator < (const edge &that) const {
        //Si se desea encontrar el árbol de recubrimiento de máxima suma, cambiar el < por
        un >
        return weight < that.weight;
    }
};

/* Union find */
int p[10001], rank[10001];
void make_set(int x){ p[x] = x, rank[x] = 0; }
void link(int x, int y){ rank[x] > rank[y] ? p[y] = x : p[x] = y, rank[x] == rank[y] ?
rank[y]++: 0; }
int find_set(int x){ return x != p[x] ? p[x] = find_set(p[x]) : p[x]; }
void merge(int x, int y){ link(find_set(x), find_set(y)); }
/* End union find */

int main(){
    int c;
    cin >> c;
    while (c--){
        int n, m;
        cin >> n >> m;
        vector<edge> e;
        long long total = 0;
        while (m--){
            edge t;
            cin >> t.start >> t.end >> t.weight;
            e.push_back(t);
            total += t.weight;
        }
    }
}

```

```

    sort(e.begin(), e.end());
    for (int i=0; i<=n; ++i){
        make_set(i);
    }
    for (int i=0; i<e.size(); ++i){
        int u = e[i].start, v = e[i].end, w = e[i].weight;
        if (find_set(u) != find_set(v)){
            //xprintf("Joining%d with%d using weight%d\n", u, v, w);
            total -= w;
            merge(u, v);
        }
    }
    cout << total << endl;
}
return 0;
}

```

2.4. Algoritmo de Floyd

```

#include <iostream>
#include <climits>
#include <algorithm>

using namespace std;

unsigned long long g[101][101];

int main(){
    int casos;
    cin >> casos;
    bool first = true;
    while (casos--){
        if (!first) cout << endl;
        first = false;

        int n, e, t;
        cin >> n >> e >> t;
        for (int i=0; i<n; ++i){
            for (int j=0; j<n; ++j){
                g[i][j] = INT_MAX;
            }
            g[i][i] = 0;
        }

        int m;
        cin >> m;
        while (m--){
            int i, j, k;
            cin >> i >> j >> k;
            g[i-1][j-1] = k;
        }

        for (int k=0; k<n; ++k){
            for (int i=0; i<n; ++i){
                for (int j=0; j<n; ++j){
                    g[i][j] = min(g[i][j], g[i][k] + g[k][j]);
                }
            }
        }
    }
}

```

```

    }
}

int r=0;
e -= 1;
for (int i=0; i<n; ++i){
    r += ((g[i][e] <= t) ? 1 : 0);
}

cout << r << endl;
}
return 0;
}

```

2.5. Puntos de articulación

```

#include <vector>
#include <set>
#include <map>
#include <algorithm>
#include <iostream>
#include <iterator>

using namespace std;

typedef string node;
typedef map<node, vector<node> > graph;
typedef char color;

const color WHITE = 0, GRAY = 1, BLACK = 2;

graph g;
map<node, color> colors;
map<node, int> d, low;

set<node> cameras;

int timeCount;

void dfs(node v, bool isRoot = true){
    colors[v] = GRAY;
    d[v] = low[v] = ++timeCount;
    vector<node> neighbors = g[v];
    int count = 0;
    for (int i=0; i<neighbors.size(); ++i){
        if (colors[neighbors[i]] == WHITE){ // (v, neighbors[i]) is a tree edge
            dfs(neighbors[i], false);
            if (!isRoot && low[neighbors[i]] >= d[v]){
                cameras.insert(v);
            }
            low[v] = min(low[v], low[neighbors[i]]);
            ++count;
        }else{ // (v, neighbors[i]) is a back edge
            low[v] = min(low[v], d[neighbors[i]]);
        }
    }
}
}

```

```

    if (isRoot && count > 1){ //Is root and has two neighbors in the DFS-tree
        cameras.insert(v);
    }
    colors[v] = BLACK;
}

int main(){
    int n;
    int map = 1;
    while (cin >> n && n > 0){
        if (map > 1) cout << endl;
        g.clear();
        colors.clear();
        d.clear();
        low.clear();
        timeCount = 0;
        while (n--){
            node v;
            cin >> v;
            colors[v] = WHITE;
            g[v] = vector<node>();
        }

        cin >> n;
        while (n--){
            node v,u;
            cin >> v >> u;
            g[v].push_back(u);
            g[u].push_back(v);
        }

        cameras.clear();

        for (graph::iterator i = g.begin(); i != g.end(); ++i){
            if (colors[(*i).first] == WHITE){
                dfs((*i).first);
            }
        }

        cout << "City map #" << map << ": " << cameras.size() << " camera(s) found" <<
endl;
        copy(cameras.begin(), cameras.end(), ostream_iterator<node>(cout, "\n"));
        ++map;
    }
    return 0;
}

```

2.6. Máximo flujo: Método de Ford-Fulkerson, algoritmo de Edmonds-Karp

El algoritmo de Edmonds-Karp es una modificación al método de Ford-Fulkerson. Este último utiliza DFS para hallar un camino de aumentación, pero la sugerencia de Edmonds-Karp es utilizar BFS que lo hace más eficiente en algunos grafos.

```

int cap[MAXN+1][MAXN+1], flow[MAXN+1][MAXN+1], prev[MAXN+1];

int fordFulkerson(int n, int s, int t){
    int ans = 0;

```



```

for (int i=0; i<n; ++i) fill(flow[i], flow[i]+n, 0);
while (true){
    fill(prev, prev+n, -1);
    queue<int> q;
    q.push(s);
    while (q.size() && prev[t] == -1){
        int u = q.front();
        q.pop();
        for (int v = 0; v<n; ++v)
            if (v != s && prev[v] == -1 && cap[u][v] > 0 && cap[u][v] - flow[u][v] > 0)
                prev[v] = u, q.push(v);
    }

    if (prev[t] == -1) break;

    int bottleneck = INT_MAX;
    for (int v = t, u = prev[v]; u != -1; v = u, u = prev[v]){
        bottleneck = min(bottleneck, cap[u][v] - flow[u][v]);
    }
    for (int v = t, u = prev[v]; u != -1; v = u, u = prev[v]){
        flow[u][v] += bottleneck;
        flow[v][u] = -flow[u][v];
    }
    ans += bottleneck;
}
return ans;
}

```

3. Programación dinámica

3.1. Longest common subsequence

```

#define MAX(a,b) ((a>b)?(a):(b))
int dp[1001][1001];

int lcs(const string &s, const string &t){
    int m = s.size(), n = t.size();
    if (m == 0 || n == 0) return 0;
    for (int i=0; i<=m; ++i)
        dp[i][0] = 0;
    for (int j=1; j<=n; ++j)
        dp[0][j] = 0;
    for (int i=0; i<m; ++i)
        for (int j=0; j<n; ++j)
            if (s[i] == t[j])
                dp[i+1][j+1] = dp[i][j]+1;
            else
                dp[i+1][j+1] = MAX(dp[i+1][j], dp[i][j+1]);
    return dp[m][n];
}

```

4. Geometría

4.1. Área de un polígono

Si P es un polígono simple (no se intersecta a sí mismo) su área está dada por:

$$A(P) = \frac{1}{2} \sum_{i=0}^{n-1} (x_i \cdot y_{i+1} - x_{i+1} \cdot y_i)$$

4.2. Centro de masa de un polígono

Si P es un polígono simple (no se intersecta a sí mismo) su centro de masa está dado por:

$$\bar{C}_x = \frac{\iint_R x dA}{M} = \frac{1}{6M} \sum_{i=1}^n (y_{i+1} - y_i)(x_{i+1}^2 + x_{i+1} \cdot x_i + x_i^2)$$

$$\bar{C}_y = \frac{\iint_R y dA}{M} = \frac{1}{6M} \sum_{i=1}^n (x_i - x_{i+1})(y_{i+1}^2 + y_{i+1} \cdot y_i + y_i^2)$$

Donde M es el área del polígono.

Otra posible fórmula equivalente:

$$\bar{C}_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i \cdot y_{i+1} - x_{i+1} \cdot y_i)$$

$$\bar{C}_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_{i+1})(x_i \cdot y_{i+1} - x_{i+1} \cdot y_i)$$

4.3. Convex hull: Graham Scan

Complejidad: $O(n \log_2 n)$

```
#include <iostream>
```

```
#include <vector>
```

```
#include <algorithm>
```

```
#include <iterator>
```

```
#include <cmath>
```

```
using namespace std;
```

```
struct point{
    int x,y;
    point() {}
    point(int X, int Y) : x(X), y(Y) {}
};
```

```
point pivot;
```

```
ostream& operator<< (ostream& out, const point& c)
{
    out << "(" << c.x << "," << c.y << ")";
    return out;
}
```

```
//P es un polígono ordenado anticlockwise.
//Si es clockwise, retorna el area negativa.
//Si no esta ordenado retorna pura mierda
```

```
double area(const vector<point> &p){
    double r = 0.0;
    for (int i=0; i<p.size(); ++i){
        int j = (i+1) % p.size();
        r += p[i].x*p[j].y - p[j].x*p[i].y;
    }
    return r/2.0;
}
```

```

}

//retorna si c esta a la izquierda de el segmento AB
inline int cross(const point &a, const point &b, const point &c){
    return (b.x-a.x)*(c.y-a.y) - (c.x-a.x)*(b.y-a.y);
}

//Self < that si esta a la derecha del segmento Pivot-That
bool angleCmp(const point &self, const point &that){
    return( cross(pivot, that, self) < 0 );
}

inline int distsqr(const point &a, const point &b){
    return (a.x - b.x)*(a.x - b.x) + (a.y - b.y)*(a.y - b.y);
}

//vector p tiene los puntos ordenados anticlockwise
vector<point> graham(vector<point> p){
    pivot = p[0];
    sort(p.begin(), p.end(), angleCmp);
    //Ordenar por ángulo y eliminar repetidos.
    //Si varios puntos tienen el mismo angulo se borran todos excepto el que esté más lejos
    for (int i=1; i<p.size()-1; ++i){
        if (cross(p[0], p[i], p[i+1]) == 0){ //Si son colineales...
            if (distsqr(p[0], p[i]) < distsqr(p[0], p[i+1])){ //Borrar el mas cercano
                p.erase(p.begin() + i);
            }else{
                p.erase(p.begin() + i + 1);
            }
        }
        i--;
    }
}

vector<point> chull(p.begin(), p.begin()+3);

//Ahora sí!!!
for (int i=3; i<p.size(); ++i){
    while ( chull.size() >= 2 && cross(chull[chull.size()-2], chull[chull.size()-1],
p[i]) < 0){
        chull.erase(chull.end() - 1);
    }
    chull.push_back(p[i]);
}

return chull;
}

int main(){
    int n;
    int tileNo = 1;
    while (cin >> n && n){
        vector<point> p;
        point min(10000, 10000);
        int minPos;
        for (int i=0; i<n; ++i){
            int x, y;
            cin >> x >> y;

```

```

    p.push_back(point(x,y));
    if (y < min.y || (y == min.y && x < min.x )){
        min = point(x,y);
        minPos = i;
    }
}
double tileArea = fabs(area(p));

//Destruye el orden cw|ccw poligono, pero hay que hacerlo por que Graham necesita el
pivot en p[0]
swap(p[0], p[minPos]);
pivot = p[0];
double hullArea = fabs(area(gham(p)));

printf("Tile #%d\n", tileNo++);
printf("Wasted Space =%.2f \\\n\n", (hullArea - tileArea) * 100.0 / hullArea);

}
return 0;
}

```

4.4. Convex hull: Andrew's monotone chain

```

Complejidad:  $O(n \log_2 n)$ 
// Implementation of Monotone Chain Convex Hull algorithm.
#include <algorithm>
#include <vector>
using namespace std;

typedef long long CoordType;

struct Point {
    CoordType x, y;

    bool operator <(const Point &p) const {
        return x < p.x || (x == p.x && y < p.y);
    }
};

// 2D cross product.
// Return a positive value, if OAB makes a counter-clockwise turn,
// negative for clockwise turn, and zero if the points are collinear.
CoordType cross(const Point &O, const Point &A, const Point &B)
{
    return (A.x - O.x) * (B.y - O.y) - (A.y - O.y) * (B.x - O.x);
}

// Returns a list of points on the convex hull in counter-clockwise order.
// Note: the last point in the returned list is the same as the first one.
vector<Point> convexHull(vector<Point> P)
{
    int n = P.size(), k = 0;
    vector<Point> H(2*n);

    // Sort points lexicographically
    sort(P.begin(), P.end());

```

```

// Build lower hull
for (int i = 0; i < n; i++) {
    while (k >= 2 && cross(H[k-2], H[k-1], P[i]) <= 0) k--;
    H[k++] = P[i];
}

// Build upper hull
for (int i = n-2, t = k+1; i >= 0; i--) {
    while (k >= t && cross(H[k-2], H[k-1], P[i]) <= 0) k--;
    H[k++] = P[i];
}

H.resize(k);
return H;
}

```

4.5. Mínima distancia entre un punto y un segmento

```

struct point{
    double x,y;
};

inline double dist(const point &a, const point &b){
    return sqrt((a.x-b.x)*(a.x-b.x) + (a.y-b.y)*(a.y-b.y));
}

inline double distsqr(const point &a, const point &b){
    return (a.x-b.x)*(a.x-b.x) + (a.y-b.y)*(a.y-b.y);
}

/*
Returns the closest distance between point pnt and the segment that goes from point a
to b
Idea by: http://local.wasp.uwa.edu.au/~pbourke/geometry/pointline/
*/
double distance_point_to_segment(const point &a, const point &b, const point &pnt){
    double u = ((pnt.x - a.x)*(b.x - a.x) + (pnt.y - a.y)*(b.y - a.y)) / distsqr(a, b);
    point intersection;
    intersection.x = a.x + u*(b.x - a.x);
    intersection.y = a.y + u*(b.y - a.y);
    if (u < 0.0 || u > 1.0){
        return min(dist(a, pnt), dist(b, pnt));
    }
    return dist(pnt, intersection);
}

```

4.6. Mínima distancia entre un punto y una recta

```

/*
Returns the closest distance between point pnt and the line that passes through points
a and b
Idea by: http://local.wasp.uwa.edu.au/~pbourke/geometry/pointline/
*/
double distance_point_to_line(const point &a, const point &b, const point &pnt){
    double u = ((pnt.x - a.x)*(b.x - a.x) + (pnt.y - a.y)*(b.y - a.y)) / distsqr(a, b);
    point intersection;
    intersection.x = a.x + u*(b.x - a.x);

```

```

    intersection.y = a.y + u*(b.y - a.y);
    return dist(pnt, intersection);
}

```

5. Java

5.1. Entrada desde entrada estándar

Este primer método es muy fácil pero es mucho más ineficiente porque utiliza Scanner en vez de BufferedReader:

```

import java.io.*;
import java.util.*;

class Main{
    public static void main(String[] args){
        Scanner sc = new Scanner(System.in);
        while (sc.hasNextLine()){
            String s= sc.nextLine();
            System.out.println("Leí: " + s);
        }
    }
}

```

Este segundo es más rápido:

```

import java.util.*;
import java.io.*;
import java.math.*;

class Main {
    public static void main(String[] args) throws IOException {
        BufferedReader reader = new BufferedReader(new InputStreamReader(System.in));
        String line = reader.readLine();
        StringTokenizer tokenizer = new StringTokenizer(line);
        int N = Integer.valueOf(tokenizer.nextToken());
        while (N-- > 0){
            String a, b;
            a = reader.readLine();
            b = reader.readLine();

            int A = a.length(), B = b.length();
            if (B > A){
                System.out.println("0");
            }else{
                BigInteger dp[][] = new BigInteger[2][A];
                /*
                dp[i][j] = cantidad de maneras diferentes
                en que puedo distribuir las primeras i
                letras de la subsecuencia (b) terminando
                en la letra j de la secuencia original (a)
                */

                if (a.charAt(0) == b.charAt(0)){
                    dp[0][0] = BigInteger.ONE;
                }else{
                    dp[0][0] = BigInteger.ZERO;
                }
                for (int j=1; j<A; ++j){

```

```

        dp[0][j] = dp[0][j-1];
        if (a.charAt(j) == b.charAt(0)){
            dp[0][j] = dp[0][j].add(BigInteger.ONE);
        }
    }

    for (int i=1; i<B; ++i){
        dp[i%2][0] = BigInteger.ZERO;
        for (int j=1; j<A; ++j){
            dp[i%2][j] = dp[i%2][j-1];
            if (a.charAt(j) == b.charAt(i)){
                dp[i%2][j] = dp[i%2][j].add(dp[(i+1)%2][j-1]);
            }
        }
    }
    System.out.println(dp[(B-1)%2][A-1].toString());
}
}
}
}
}
}
}
}
}
}

```

5.2. Entrada desde archivo

```

import java.io.*;
import java.util.*;
public class BooleanTree {
    public static void main(String[] args) throws FileNotFoundException {
        System.setIn(new FileInputStream("tree.in"));
        System.setOut(new PrintStream("tree.out"));
        Scanner reader = new Scanner(System.in);
        N = reader.nextInt();
        for (int c = 1; c <= N; ++c) {
            int res = 100;
            if (res < 1000)
                System.out.println("Case #" + c + ": " + res);
            else
                System.out.println("Case #" + c + ": IMPOSSIBLE");
        }
    }
}

```

6. C++

6.1. Entrada desde archivo

```

#include <iostream>
#include <fstream>

using namespace std;

int _main(){
    freopen("entrada.in", "r", stdin);
    freopen("entrada.out", "w", stdout);

    string s;
    while (cin >> s){
        cout << "Leí " << s << endl;
    }
}

```

```

    }
    return 0;
}

int main(){
    ifstream fin("entrada.in");
    ofstream fout("entrada.out");

    string s;
    while (fin >> s){
        fout << "Leí " << s << endl;
    }
    return 0;
}

```

6.2. Strings con caracteres especiales

```

#include <iostream>
#include <cassert>
#include <stdio.h>
#include <assert.h>
#include <wchar.h>
#include <wctype.h>
#include <locale.h>

using namespace std;

int main(){
    assert(setlocale(LC_ALL, "en_US.UTF-8") != NULL);
    wchar_t c;

    wstring s;
    while (getline(wcin, s)){
        wcout << L"Leí : " << s << endl;
        for (int i=0; i<s.size(); ++i){
            c = s[i];
            wprintf(L"%lc%lc\n", tolower(s[i]), toupper(s[i]));
        }
    }

    return 0;
}

```

Nota: Como alternativa a la función `getline`, se pueden utilizar las funciones `fgetws` y `fputws`, y más adelante `swscanf` y `wprintf`:

```

#include <iostream>
#include <cassert>
#include <stdio.h>
#include <assert.h>
#include <wchar.h>
#include <wctype.h>
#include <locale.h>

using namespace std;

```



```

int main(){
    assert(setlocale(LC_ALL, "en_US.UTF-8") != NULL);
    wchar_t in_buf[512], out_buf[512];
    swprintf(out_buf, 512, L"¿Podrías escribir un número?, Por ejemplo%d. ¡Gracias,
pingüino español!\n", 3);
    fputws(out_buf, stdout);

    fgetws(in_buf, 512, stdin);
    int n;
    swscanf(in_buf, L"%d", &n);

    swprintf(out_buf, 512, L"Escribiste%d, yo escribo ¿ÏàÛÑ~\n", n);
    fputws(out_buf, stdout);

    return 0;
}

```